A New Class of Wilf-Equivalent Permutations

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Abstract. For about 10 years, the classification up to Wilf equivalence of permutation patterns was thought completed up to length 6. In this paper, we establish a new class of Wilf-equivalent permutation patterns, namely, $(n-1, n-2, n, \tau) \sim (n-2, n, n-1, \tau)$ for any $\tau \in S_{n-3}$. In particular, at level n = 6, this result includes the only missing equivalence (546213) ~ (465213), and for n = 7 it completes the classification of permutation patterns by settling all remaining cases in S_7 .

Keywords: Wilf-equivalent, shape-Wilf-equivalent, restricted patterns, forbidden permutations

1. Introduction

A permutation τ of length k is written as $(a_1, a_2, ..., a_k)$ where $\tau(i) = a_i, 1 \le i \le k$. For k < 10 we suppress the commas without causing confusion. As usual, S_n denotes the symmetric group on $[n] = \{1, 2, ..., n\}$.

Definition 1 Let τ and π be two permutations of lengths k and n, respectively. We say that π is τ -avoiding if there is no subsequence $i_{\tau(1)}, i_{\tau(2)}, \ldots, i_{\tau(k)}$ of [n] such that $\pi(i_1) < \pi(i_2) < \cdots < \pi(i_k)$. If there is such a subsequence, we say that the subsequence $\pi(i_{\tau(1)}), \pi(i_{\tau(2)}), \ldots, \pi(i_{\tau(k)})$ is of type τ .

For example, the permutation $\omega = (52687431)$ avoids (2413) but does not avoid (3142) because of its subsequence (5283). An equivalent, but perhaps more insightful, definition is the following reformulation in terms of matrices.

Definition 2 Let $\tau \in S_n$. The *permutation matrix* $M(\tau)$ is the $n \times n$ matrix having a 1 in position $(i, \tau(i))$ for $1 \le i \le n$, and having 0 elsewhere.¹ Given two permutation matrices M and N, we say that M avoids N if no submatrix of M is identical to N.

Note that a permutation matrix M of size n is simply a traversal of an $n \times n$ matrix, i.e. an arrangement of 1's for which there is exactly one 1 in every row and in every column of M. It is clear that a permutation $\pi \in S_n$ contains a subsequence $\tau \in S_k$ if and only if $M(\pi)$ contains $M(\tau)$ as a submatrix. Let $S_n(\tau)$ denote the set of τ -avoiding permutations in S_n .

Definition 3 Two permutations τ and σ are called *Wilf-equivalent* if they are equally restrictive: $|S_n(\tau)| = |S_n(\sigma)|$ for all $n \in \mathbb{N}$. We denote this by $\tau \sim \sigma$. If $|S_k(\tau)| = |S_k(\sigma)|$ for $k \leq n$, then we say that τ and σ are *equinumerant* up to level n.

The basic problem in the theory of forbidden subsequences is to classify all permutations up to Wilf-equivalence. Obviously, if two permutations are Wilf-equivalent, then they must be of the same length. Further, many Wilf-equivalences can be deduced by symmetry arguments within the same S_k . For instance, if $M(\pi)$ contains $M(\tau)$ as a submatrix, then the transpose matrix $M(\pi)^t$ contains $M(\tau)^t$. The same is true when simultaneously reflecting both matrices $M(\pi)$ and $M(\tau)$ in either a horizontal or a vertical axis of symmetry. The three operations defined above generate the dihedral group D_4 acting on the set of permutation matrices in the obvious way. The orbits of D_4 in S_k are called symmetry classes. It is clear that if τ and σ belong to the same symmetry class in S_k , then $\tau \sim \sigma$. However, Wilf-classes are in general, but apparently rarely, larger than single symmetry classes. This makes the classification of permutations up to Wilf-equivalence a subtle and difficult process.

The first major result in the theory of forbidden subsequences states that $(123) \sim (132)$, and hence S_3 is one Wilf-class, which combines the two symmetry classes of (123) and (132). At the behest of Wilf, bijections between $S_n(123)$ and $S_n(132)$ were given by Simion-Schmidt [15], Rotem [14], Richards [13], and West [19]. They all prove $|S_n(123)| = c_n$, where c_n is the *n*th Catalan number. Permutations with forbidden subsequences arise naturally in computer science in connection with sorting problems and strings with forbidden subwords. For example, in [6–7] Knuth shows that $S_n(231)$ is the set of stack-sortable permutations (see also [9]), so that $|S_n(231)|$ is the number of binary strings of length 2n, in which 0 stands for a "move into a stack" and 1 symbolizes a "move out from the stack".

Numerous problems involving forbidden subsequences have also appeared in algebraic combinatorics. In the late 1980s, it was discovered that the property of avoiding 2143 exactly characterizes the *vexillary* permutations, i.e. those whose Stanley symmetric function is a Schur function. (See [10] for a good exposition.) Lakshimibai and Sandhya [8] likewise show that $S_n(3412, 4231)$ is the set of permutations indexing an interesting subclass of Schubert varieties. And Billey and Warrington [3] have very recently defined a class of permutations under 5 restrictions which are related to the Kazhdan-Lusztig polynomials. This all naturally leads to the study and classification of Wilf-classes of permutations of length 4 or more.

The classification of S_4 turns out to be much more complicated than that of S_3 . It is completed in a series of papers by Stankova and West. They utilize the concept of a *generating* tree $\mathfrak{T}(\tau)$ of $\tau \in S_k$: the nodes on the *k*th level of $\mathfrak{T}(\tau)$ are the permutations in $S_n(\tau)$, and the descendants of $\pi \in S_n(\tau)$ are obtained from π by inserting n + 1 in appropriate places in π so that τ is still avoided. Clearly, the tree isomorphism $\mathfrak{T}(\tau) \simeq \mathfrak{T}(\sigma)$ implies $\tau \sim \sigma$, but the converse is far from true. In [19], West shows $\mathfrak{T}(1234) \simeq \mathfrak{T}(1243) \simeq \mathfrak{T}(2143)$. In [16], Stankova constructs a specific isomorphism $\mathfrak{T}(4132) \cong \mathfrak{T}(3142)$. In [17], she completes the classification of S_4 by proving (1234) ~ (4123); there she uses a different approach

$\tau \in S_4$	$S_6(\tau)$	$S_7(\tau)$	$S_8(\tau)$	$S_9(\tau)$	$S_{10}(\tau)$	$S_{11}(\tau)$	$S_{12}(\tau)$	$S_{13}(au)$
(4132)	512	2740	15485	91245	555662	3475090	22214707	144640291
(1234)	513	2761	15767	94359	586590	3763290	24792705	167078577
(1324)	513	2762	15793	94776	591950	3824112	25431452	173453058

Figure 1. Classification of S₄ up to Wilf-equivalence.



Figure 2. Splitting of a traversal $T \in S_{Y(10,10,9,8,8,8,8,7,4,3)}(213)$.

which yields the somewhat surprising result that, while $\mathfrak{T}(1234) \not\cong \mathfrak{T}(4123)$, on every level of the two trees the number of nodes with a given number of descendants is the same for both trees. Thus, the seven symmetry classes of S_4 are grouped in three Wilf-classes, with representatives (4132), (1234) and (1324) (cf. figure 1.)

In [1], Babson-West show $(n-1, n, \tau) \sim (n, n-1, \tau)$ for any $\tau \in S_{n-2}$, and $(n-2, n-1, n, \tau) \sim (n, n-1, n-2, \tau)$ for any $\tau \in S_{n-3}$, thus completing the classification up to level 5. The key idea is the concept of a stronger Wilf-equivalence relation.

Definition 4 A *traversal* T of a Young diagram Y is an arrangement of 1's and 0's such that every row and every column of Y has exactly one 1 in it. A subset of 1's in T is said to form a *submatrix* of Y if all columns and rows of Y passing through these 1's intersect *inside* Y. For a permutation $\tau \in S_k$, we say that T contains pattern τ if some k 1's of T form a submatrix of Y identical to $M(\tau)$ (cf. figure 2.)

Given several 1's in a traversal *T*, the condition for them to form a submatrix of *T* is the same as the requirement that the column of the rightmost 1 and the row of the lowest 1 must intersect *inside Y*. This condition is necessary for the new definition to be a useful generalization of the classical definition of a forbidden subsequence, as we shall see below. In particular, when *Y* is a square diagram, the two definitions coincide. Let us denote by $S_Y(\tau)$ the set of all traversals of *Y* which avoid τ .

Definition 5 Two permutations τ and σ are called *shape-Wilf-equivalent* (SWE) if $|S_Y(\tau)| = |S_Y(\sigma)|$ for all Young diagrams *Y*. We denote this by $\tau \stackrel{s}{\sim} \sigma$.

Clearly, $\tau \sim \sigma$ implies $\tau \sim \sigma$, but not conversely. We will write $Y(a_1, a_2, ..., a_n)$ for the Young diagram Y whose *i*-th row has a_i cells, $1 \le i \le n$. In order for a Young diagram Y to have any traversals at all, Y must have the same number of rows and columns and Y must contain the *staircase* diagram St = Y(n, n - 1, ..., 2, 1), where n is the number

of cells in the top (largest) row of Y. Thus, from now on, when we talk about a Young diagram Y of size n, we will assume that Y has n rows and n columns and contains St of size n.

SWE is a very strong relation on two permutations and it is certainly too restrictive on its own to be useful in the general classification of permutations. However, combined with the proposition below (see [1]), it allows for more Wilf-equivalences to be established.

Proposition 1 Let $A \stackrel{s}{\sim} B$ for some permutation matrices A and B. Then for any permutation matrix C:

$$\left(\begin{array}{c|c} A & 0\\ \hline 0 & C \end{array}\right) \stackrel{s}{\sim} \left(\begin{array}{c|c} B & 0\\ \hline 0 & C \end{array}\right).$$

Let I_k be the $k \times k$ identity matrix, and let J_k be its reflection across a vertical axis of symmetry. According to Backelin-West-Xin in [2], $I_k \stackrel{s}{\sim} J_k$ for any k, and hence $(n, n - 1, ..., m, \tau) \sim (m, ..., n - 1, n, \tau)$ for any $\tau \in S_{n-m}$. This SWE generalizes the results in [20] and [1], but it is not sufficient to complete the classification of S_7 , nor of S_6 .

In 2001, Stankova noticed a missing case of a plausible Wilf-equivalence in S_6 : (546213) and (465213) were equinumerant up to level 11, but no reference was found regarding why these permutations were thought to be in different Wilf classes (see figure 11). Stankova further found an infinite class of Wilf-equivalences

$$(n-1, n-2, n, \tau) \sim (n-2, n, n-1, \tau),$$
 (1)

At her request, West confirmed (1) by computer checks for n = 6, 7 up to level 13.²

The purpose of this paper is to explain the proof of the new Wilf-equivalences in (1). The idea is to show $(213) \stackrel{s}{\sim} (132)$ and apply then Proposition 1. Even though M(213) and M(132) are transposes of each other, their SWE relationship is far from trivial, as the present paper will reveal. It is surprising that such a basic relationship is discovered only now, 10 years after the introduction of SWE in the early 1990s.

Theorem 1 (*Main result of the paper*) The permutations (213) and (132) are shape-Wilf-equivalent. Consequently, for any $\tau \in S_{n-3}$, the permutations $(n - 1, n - 2, n, \tau)$ and $(n - 2, n, n - 1, \tau)$ are Wilf-equivalent.

Theorem 1 finally accounts for the last missing case in S_6 and the remaining cases in S_7 , thus completing the classification of forbidden subsequences up to length n = 7 (see figure 7).

In summary, modulo symmetry classes, as of now there are essentially two known infinite families of Wilf-equivalences, resulting from [2] and the present paper:

$$\left(\frac{I_k \mid 0}{0 \mid C}\right) \stackrel{s}{\sim} \left(\frac{J_k \mid 0}{0 \mid C}\right) \quad \text{and} \quad \left(\frac{M(213) \mid 0}{0 \mid C}\right) \stackrel{s}{\sim} \left(\frac{M(132) \mid 0}{0 \mid C}\right). \tag{2}$$

Further, there is only one known "sporadic" case of Wilf-equivalence, from [16]:

$$\begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 \end{pmatrix} \sim \begin{pmatrix} 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 \end{pmatrix}.$$
 (3)

The above (4132) and (3142) in (3) constitute an interesting pair of Wilf-equivalent permutations: (3142) has the smallest symmetry class as it corresponds geometrically to the quadrilateral with most symmetries—the square, while (4132) has the largest symmetry class as it corresponds to the quadrilateral with least number of symmetries—a quadrilateral with 4 different angles. And yet, not only (4132) \sim (3142), but also their trees are isomorphic. In a similar vein, the permutations in (2) are more than just Wilf-equivalent they are SWE. This is an interesting phenomenon—so far, every known Wilf-equivalence can be explained by a stronger relationship: either symmetry, tree isomorphism, SWE, or a combination of these.

For further discussion, we refer the reader to Sections 5–6.

The proof of the main result (Theorem 1) is structured as follows. Since the permutation matrices of (213) and (132) are transposes, Lemma 2(i)–(ii) in Section 3 allows us to prove the equivalent statement that $Y \equiv Y^t$ for all Young diagrams Y, i.e. that the number of (213)-avoiding traversals of Y is the same as the number of (213)-avoiding traversals of Y is the same as the number of (213)-avoiding traversals of Y is the same as the number of (213)-avoiding traversals of Y^t (cf. Definition 10 in Section 2). To this end, in Section 2 we define an operation on Young diagrams, which we call row-decomposition. It breaks up every Young diagram Y into two smaller diagrams Y'_r and Y''_r so that $Y \equiv Y'_r + Y''_r$. To link this with the transpose Y^t , we define an analogous column-decomposition on Y, and Lemma 2(iii) shows that establishing $Y \equiv Y'_c + Y''_c$ is equivalent to the main result (213) \sim (132). In Section 4, we investigate a commutativity argument: we apply row decomposition followed by column decomposition, then reverse the order of the two operations and compare the results. This gives us a tool for using induction and the row-decomposition formula to prove the desired column decomposition formula. An amusing consequence of this discussion in presented by Corollary 2 in Section 5.

2. The row-decomposition formula

Let *Y* be a Young diagram with *n* rows and *n* columns. We denote by (k, ℓ) the intersection cell of row *k* and column ℓ , counted from the top left corner of *Y*. A cell in the bottom row of *Y* is called a *bottom* cell of *Y*. Let *m* be the number of bottom cells in *Y*.

Definition 6 For a subset X of cells in Y, define the *reduction* $Y/_X$ of Y along X to be the new Young diagram obtained from Y by deleting all rows and columns of Y which intersect X.

For example, if the *cross* of a cell C in Y is the union of the row and column containing C, then the reduction $Y/_C$ is the diagram obtained from Y by deleting the cross of C. The



Figure 3. $Y \rightarrow Y_C \rightarrow Y_C = A_C \times B_C$.

reduction of Y along an arbitrary bottom cell of Y is denoted by Y^r , and it clearly does not depend on the choice of the bottom cell. We will use this fact frequently when reducing along cell (1, n) (and (n, 1)) in the proof of the commutativity argument in Lemma 3 in Section 4.

Definition 7 To any bottom cell *C* in *Y*, we associate a *cross-product* Y_C of Young diagrams in the following way. Mark by *P* the top right corner of *C* (this is a grid point on *Y*), and consider the reduction Y_C , which still contains point *P*. Starting from *P* in Y_C , draw a 45° ray in north-east direction until the ray intersects for the first time the border of Y_C , and use the resulting segment as the south-west/north-east diagonal of a smaller subdiagram A_C of Y_C . Delete the rows and columns of A_C in Y_C , leaving a subdiagram $B_C = Y_{\{C,A_C\}}$. Thus, any bottom cell *C* in the original diagram *Y* determines a pair (A_C , B_C) of smaller Young diagrams, which we call the *cross-product* of A_C and B_C and denote by $Y_C := A_C \times B_C$.

If one of the subdiagrams A_C or B_C is empty, we define Y_C to equal the other subdiagram. This case occurs exactly when C is the first or the last bottom cell of Y:

 $Y_{(n,1)} = Y^r \times \emptyset = Y^r = \emptyset \times Y^r = Y_{(n,m)}.$

Example 1 Let Y = Y(10, 10, 9, 8, 8, 8, 8, 7, 4, 3). Let C = (10, 2). Then $Y_C = A_C \times B_C = Y(6, 6, 6, 6, 5, 2) \times Y(3, 3, 2)$ (cf. figure 3.)

Definition 8 Let *Y* be a Young diagram of size *n*. The *row decomposition* of *Y* is the formal sum $\mathcal{R}(Y)$ of cross-products of smaller Young diagrams:

$$\mathcal{R}(Y) := \sum_{C} Y_{C} = \sum_{C} A_{C} \times B_{C},$$

where the sum is taken over all bottom cells C of Y.

As noted above, the first and the last summands of $\mathcal{R}(Y)$ are identical to Y^r .

Definition 9 A traversal of a diagram *Y* which avoids (resp. contains) a (213)-pattern is called a *good* (resp. *bad*) traversal of *Y*. Denote by $\mathcal{T}(Y) = \mathcal{T}(a_1, a_2, ..., a_n)$ the number of good traversals of $Y = Y(a_1, a_2, ..., a_n)$.

We use the convention $\mathcal{T}(\emptyset) = 1$. For some 1 already placed in a cell of *Y*, we say that it *imposes a* (213)-*condition on Y* if it plays the role of a "1" in a (213)-pattern contained in some bad traversal of *Y*; the (213)-condition is the actual condition on the rest of *Y* in order to avoid a (213)-pattern containing this 1.

Definition 10 Two diagrams *Y* and *X* are said to be *numerically equivalent* if $\mathcal{T}(Y) = \mathcal{T}(X)$. We denote this by $Y \equiv X$.

Clearly, $\mathcal{T}(A_C \times B_C) = \mathcal{T}(A_C) \cdot \mathcal{T}(B_C)$. Moreover, to obtain the number $\mathcal{T}(Y)$, we can apply the function \mathcal{T} to all terms in the formal sum $\mathcal{R}(Y)$:

Theorem 2 (*Row-Decomposition*) Let Y be a Young diagram of size n. Then

$$\mathcal{T}(Y) = \sum_{C} \mathcal{T}(A_{C}) \cdot \mathcal{T}(B_{C})$$
(4)

where the sum is taken over all bottom cells C in Y.

Proof: For a bottom cell C = (n, i), let Y_i denote the diagram Y with the additional data of 1 in cell C. The correspondence will be induced by the map $Y_i \mapsto (A_C, B_C)$. In fact, we claim that the good traversals of Y_i are in 1-1 correspondence with the good traversals of the pair of diagrams (A_C, B_C) , and hence

$$\mathcal{T}(Y_i) = \mathcal{T}(A_C) \cdot \mathcal{T}(B_C). \tag{5}$$

This, and the fact that any traversal of Y must contain exactly one 1 in the bottom row, immediately establishes Theorem 2. Hence, it suffices to prove (5) for Y_i .

When i = 1 or i = m, the claim (5) is trivial. Indeed, in these cases, the 1 in the bottom row of *Y* is either in the first or last bottom cell, and hence it doesn't impose any (213)-conditions on *Y*. The question reduces to finding all good traversals of $Y_{C} = Y^{r} = Y_{C}$. Therefore, $Y_{1} \equiv Y_{m} \equiv Y^{r}$.

Assume now that 1 < i < m. Fix a good traversal *T* of Y_i . Denote by 1_j the 1 in column *j*. By $1_j > 1_k$ we mean that 1_j is in a row above the row of 1_k . Similarly, for two disjoint sets *A* and *B* of 1's, by A > B we mean that all 1's in *A* are above all 1's in *B*. Let $B_L = \{1_1, 1_2, \ldots, 1_{i-1}\}$ denote the set of all 1's in *T* appearing in columns to the left of cell *C*. Similarly, let $A = \{1_{i+1}, 1_{i+2}, \ldots, 1_k\}$ be the set of all 1's appearing in the columns of *Y* intersecting A_C , and let $B_R = \{1_{k+1}, 1_{k+2}, \ldots, 1_n\}$ be the set of all 1's appearing in the remaining columns, i.e. all columns to the right of A_C . Notice that no 1 in B_R can appear in a row intersecting A_C via the 45° segment.

The key idea of the proof is contained in the following lemma:

Lemma 1 In any good traversal T of Y_i , we have $B_L > A$.

Proof (of Lemma 1): Given a row j of A, let the *level* L_j be the subset of A consisting of all 1's whose orthogonal projections onto row j are inside A_C . Clearly, every $1 \in A$

belongs to at least one level L_j , and $L_{n-1} \subseteq L_{n-2} \subseteq \cdots \subseteq L_{n-k} = A$, where *k* is the size of A_C . Recall that i < m, so that $L_{n-1} \cap A \neq \emptyset$, and that the bottom row of *Y* is filled by 1_i . This imposes a (213)-condition on Y_i : $B_L > L_{n-1}$. Finally, from the construction of A_C , $|L_{n-j}| > j$ for $j = 1, 2, \ldots, k-1$, and $|L_{n-k}| = k$.

We will prove simultaneously the following two statements for all $1 \le j \le k$:

- (i) $B_L > L_{n-j}$;
- (ii) All rows n 1, n 2, ..., n j of Y_i are filled in with 1's on level L_{n-j} .

For j = 1, (i) was shown above. But then row n - 1 in Y_i must be filled with an element of L_{n-1} , so (ii) is also true. Assume (i) and (ii) for some j < k. Then (ii) for j, together with $|L_{n-j}| > j$, implies that at least one element 1_s of L_{n-j} is in row n - j - 1 or above. Now (i) for j implies in particular $B_L > 1_s$, so that a new (213)-condition is imposed: $B_L > (L_{n-j-1} \setminus L_{n-j})$. Combining, $B_L > L_{n-j-1}$: this is (i). By definition of L_{n-j-1} , the only 1_h that can possibly fill in row n - j - 1 in Y_i must belong either to L_{n-j-1} , or to B_L . Because $|L_{n-j-1}| \ge j + 1$ and $B_L > L_{n-j-1}$, we conclude that $1_h \in L_{n-j-1}$. This shows (ii) for j + 1 and completes the inductive proof of the above statement. Lemma 1 follows automatically from (i) for j = k.

End of Proof of Theorem 2: Combining Lemma 1 with a previous observation, we see that no 1's from B_L or from B_R can fill the rows intersecting A_C . In other words, all rows of A_C must be filled exactly with the 1's from set A: the number of necessary 1's to make a traversal of A_C matches |A| = k because, by construction, A_C has as many rows as columns. This in turn forces all 1's in B_L and B_R to make a good traversal of the subdiagram B_C . It remains to show that there are no further (213)-conditions imposed by a triple of 1's coming from A_C and B_C .

The only way for A_C and B_C to engage together in a (213)-pattern is to have the "2" in B_L , the "1" in A, and the "3" in B_R ; or, to have the "2" and the "1" in A, and the "3" in B_R . Even though such configurations of three 1's are possible, their "full" matrices will not be contained entirely in Y because of the relative positioning of A_C and B_R .

Putting everything together, the good traversals of Y_i are in 1 - 1 correspondence with pairs of good traversals of A_C and B_C , i.e. $\mathcal{T}(Y_i) = \mathcal{T}(A_C) \cdot \mathcal{T}(B_C)$ for all bottom cells C of Y. This completes the proof of the Row-Decomposition formula.

Example 2 To illustrate the above proof, consider Y(10, 10, 9, 8, 8, 8, 7, 4, 3). Let Y_2 denote the diagram Y with the additional data that the 1 in the bottom row is in cell C = (10, 2). We have to show that $\mathcal{T}(Y_2) = \mathcal{T}(6, 6, 6, 6, 5, 2) \cdot \mathcal{T}(3, 3, 2)$ (cf. figure 3.)

The initial condition that 1_2 is in the bottom row forces the (213)-condition $1_1 > 1_3$. Since 1_3 is above row 10, the only 1's which can fill row 9 are 1_3 and 1_4 . If 1_3 is in row 9, then $1_1 > 1_4$ in order to avoid (213); if 1_4 is in row 9, then $1_1 > 1_3 > 1_4$. In any case, $1_1 > 1_3$, 1_4 . Without loss of generality, assume that 1_3 is in row 9, so that 1_4 is in row 8 or above. From $1_1 > 1_4$ and avoiding (213), we conclude that $1_1 > 1_4$, 1_5 , 1_6 , 1_7 . One of the latter four 1's must fill in row 8. Without loss of generality, assume that 1_4 is in row 8; hence 1_5 , 1_6 , 1_7 are in rows 7 or above. But then, to avoid (213), we are forced to conclude that $1_1 > 1_8$, i.e. 1_1 is above all of $1_3, \ldots, 1_8$. We need six 1's to fill in the six rows 9, $8, \ldots, 4$. It



Figure 4. $T(9, 9, 9, 9, 7, 7, 7, 7, 4) = T(9, 9, 9, 9, 7, 7, 7, 7, 3) + T(4, 4, 4, 4)^2$.

immediately follows that $1_3, \ldots, 1_8$ must have filled all 6 rows and columns of subdiagram A_C , leaving all remaining 1's (except for 1_2) to form a good traversal of subdiagram B_C (cf. figure 2.)

We note that B_C consists of two disjoint parts: $B_L(1, 1, 1)$ and $B_R(2, 2, 1)$. The argument that B_C and A_C cannot engage together in a pattern (213) is identical to the corresponding part of the proof of Theorem 2.

Given a diagram Y, let C_b be its rightmost bottom cell, which we call the *bottom corner* of Y, and let C_{b-1} be the bottom cell to the left of C_b . (In our previous notation, $C_b = (n, m)$, $C_{b-1} = (n, m-1)$.) Deleting C_b from Y results in a new diagram, which we denote by Y'_r and call the *row-deletion of* Y. Similarly, we define the *right corner* C_t of Y as the bottom cell in the rightmost column of Y, C_{t-1} as the cell directly above C_t , and the *column-deletion* Y'_c by deleting C_t from Y. Note that $(Y'_r)^t = (Y^t)'_c$, where X^t denotes as usual the transpose of diagram X along its main (north-west to south-east) diagonal.

In the row-decomposition of *Y*, we distinguish one special summand: the last but one summand $Y_{C_{b-1}} = A_{C_{b-1}} \times B_{C_{b-1}}$, which we denote by Y_r'' .

Corollary 1 For any diagram $Y, Y \equiv Y'_r + Y''_r$.

Proof: The row-decomposition of *Y* includes one more summand than the row-decomposition of Y'_r , namely, Y''_r : $\mathcal{R}(Y) = Y'_r + Y''_r$. Theorem 2 completes the proof.

From now on, we shall refer to Corollary 1 as Row-Decomposition (RD).

Example 3 Figure 4 illustrates Corollary 1 for $\mathcal{T}(9, 9, 9, 9, 7, 7, 7, 7, 4)$.

3. Column decomposition

Definition 11 The *column decomposition* C(Y) of Y is defined by:

$$\mathcal{C}(Y) = \sum_{C} ((Y^t)_{C^t})^t$$

where C runs over all cells in the rightmost column of Y, and C^t is the image of the cell C after transposing Y.

Note that the column decomposition C(Y) can be obtained directly from *Y* without going through the transpose *Y*^{*t*}: for a cell *C* in the rightmost column of *Y*, mark the south-west corner of *C*, draw the 45° ray in south-west direction until it intersects the border of *Y*, delete the cross of *C* and define analogously the product $Y_C = A_C \times B_C$.

As with row decomposition, we denote by Y'_c the diagram resulting from Y by deleting the right corner cell C_t , and by Y''_c the summand in $\mathcal{C}(Y)$ corresponding to the cell C_{t-1} right above C_t . By definition, it is clear that

$$\mathcal{C}(Y) = Y_c' + Y_c''$$

It is not obvious, however, why the same formula should be true *after* applying the function T to all terms: why is $Y \equiv Y'_c + Y''_c$?

Lemma 2 The following statements are equivalent:

(i) (213) ^s (132).
 (ii) Y ≡ Y^t for all Young diagrams Y.
 (iii) Y ≡ Y'_c + Y''_c.

Proof: Let $\mathcal{T}_{\tau}(Y)$ be the number of traversals of *Y* avoiding permutation τ . In our previous notation, $\mathcal{T}(Y) = \mathcal{T}_{(213)}(Y)$. Since $M(132) = M(213)^t$, $\mathcal{T}_{(132)}(Y) = \mathcal{T}_{(213)}(Y^t)$. By definition of SWE, (i) and (ii) are equivalent.

We will show the equivalence of (ii) and (iii) by induction on the size *n* of *Y*. When $n \le 3$, (ii) and (iii) can be easily checked by hand. Note that from the definitions of column reduction, deletion and decomposition, $(Y^t)'_c = (Y'_r)^t$ and $(Y^t)''_c = (Y''_r)^t$. Assume now that (ii) and (iii) are equivalent for size $\le n$.

Assume first that (iii) is true for size $\leq n + 1$. Then we can use (ii) for sizes $\leq n$:

$$Y^{t} \stackrel{\text{(iii)}}{\equiv} (Y^{t})_{c}^{\prime} + (Y^{t})_{c}^{\prime\prime} \stackrel{\text{def}}{=} (Y_{r}^{\prime})^{t} + (Y_{r}^{\prime\prime})^{t} \stackrel{\text{(iii)}}{\equiv} Y_{r}^{\prime} + Y_{r}^{\prime\prime} \stackrel{\text{RD}}{\equiv} Y.$$

This shows $Y \equiv Y^t$ and completes the proof of (iii) \Rightarrow (ii) for size n + 1. Conversely, assume that (ii) is true for size $\leq n + 1$.

$$Y^{t} \stackrel{\text{(ii)}}{\equiv} Y \stackrel{\text{RD}}{\equiv} Y'_{r} + Y''_{r} \stackrel{\text{(ii)}}{\equiv} (Y'_{r})^{t} + (Y''_{r})^{t} \stackrel{\text{def}}{=} (Y^{t})'_{c} + (Y^{t})''_{c}.$$
(6)

Replacing *Y* by *Y*^t in (6), the above reads $Y \equiv Y'_c + Y''_c$. This shows (ii) \Rightarrow (iii) for size n + 1, and completes the proof of (ii) \Leftrightarrow (iii) and of Lemma 2.

From now on, we shall refer to statement (iii) as Column-Decomposition (CD).

4. Commutativity of row and column decompositions

Lemma 3 $Y \equiv Y'_c + Y''_c$ for all Young diagrams Y.

Proof: Assume that the statement is true for all diagrams of size smaller than the size of Y. The idea is to apply RD and the assumed CD one after the other in different orders: this results in representing both sides of the equality as sums of the same four terms. Let us start with Y. Recall that C_b and C_t are the bottom and right corners of Y, respectively. Apply first RD:

$$Y \stackrel{\text{RD}}{\equiv} Y'_r + Y''_r.$$

Next, apply the assumed CD to Y'_r , and to the factor of Y''_r that still contains C_t , leaving the other factor of Y''_r unchanged:

$$Y'_{r} \stackrel{\text{CD}}{=} (Y'_{r})'_{c} + (Y'_{r})''_{c}, \quad Y''_{r} \stackrel{\text{CD}}{=} (Y''_{r})'_{c} + (Y''_{r})''_{c}.$$
(7)

Now start with $Y'_c + Y''_c$, and apply RD to Y'_c , and to the factor of Y''_c that still contains C_b , leaving the other factor of Y''_c unchanged:

$$Y'_{c} + Y''_{c} \stackrel{\text{RD}}{=} (Y'_{c})'_{r} + (Y'_{c})''_{r} + (Y''_{c})'_{r} + (Y''_{c})''_{r}.$$
(8)

In both Eqs. (7) and (8), by abuse of notation, we wrote $(Y''_r)'_c, (Y''_r)''_c, (Y''_c)'_r$ and $(Y''_c)''_r$ for the cross-products of diagrams resulting from applying CD, resp. RD, to the factors containing C_t , resp. C_b . From the definition of row and column deletion, it is clear that the first terms in (7) and (8) are equal: $(Y'_r)'_c = Y - C_b - C_t = (Y'_c)'_r$. We claim that the remaining three terms also pair up as:

$$(Y'_r)''_c = (Y''_c)'_r, \quad (Y''_r)'_c = (Y'_c)''_r, \quad (Y''_r)''_c = (Y''_c)''_r, \tag{9}$$

except in Case IV below where

$$(Y'_r)''_c \equiv (Y'_c)''_r, \quad (Y''_r)'_c \equiv (Y''_c)'_r, \quad (Y''_r)''_c \equiv (Y''_c)''_r.$$
(10)

Before we embark on the proofs of (9–10), note how they fit in the general outline of the proof of CD:

$$Y \stackrel{\text{RD}}{=} Y'_{r} + Y''_{r} \text{ (known)}$$

$$\stackrel{\text{CD}}{=} (Y'_{r})'_{c} + (Y'_{r})''_{c} + (Y''_{r})'_{c} + (Y''_{r})''_{c} \text{ (assumed)}$$

$$\equiv (Y'_{c})'_{r} + (Y'_{c})''_{r} + (Y''_{c})'_{r} + (Y''_{c})''_{r} \text{ (by examining cases below)}$$

$$\equiv Y'_{c} + Y''_{c} \text{ (converse RD, known).}$$

One of the reasons that this works is that the CD-factor in Y'_r can be anticipated from \mathcal{B} , the CD-factor in the original Y.

The proof of (9–10) depends solely on how the row and column decomposition interact with each other in any given Young diagram Y, more precisely, on the relative position of the two 45° segments used in the decompositions. Let the *RD-segment* be the segment used in RD, and let the *RD-factor* be the subdiagram A_C determined by the RD-segment, and similarly for CD. Set $\mathcal{A} := \text{RD-factor}$, and $\mathcal{B} := \text{CD-factor}$ in Y. To see that Eqs. (9) and (10) are true, divide all Young diagrams Y into four cases: since the RD- and CD-segments are parallel to each other, there are only four possible relative positions for them. We will use $\stackrel{\text{sym}}{=}$ as shorthand for "by symmetry arguments".

Case I The RD- and CD-factors do not overlap $(A \cap B = \emptyset)$, i.e. the RD- and CD-segments hit *Y*'s border before they "come close" to each other (figure 12). Then

$$\begin{aligned} (Y'_r)''_c &= (Y - C_b)''_c = (Y - C_b)/_{\{(1,n),\mathcal{B}\}} \times \mathcal{B} \\ &= (Y/_{\{(1,n),\mathcal{B}\}})'_r \times \mathcal{B} = (Y''_c)'_r; \\ \stackrel{\text{sym}}{\Rightarrow} (Y'_c)''_r &= (Y - C_t)/_{\{(n,1),\mathcal{A}\}} \times \mathcal{A} = (Y''_r)'_c; \\ (Y''_r)''_c &= (Y/_{\{(n,1),\mathcal{A}\}} \times \mathcal{A})''_c = (Y/_{\{(n,1),\mathcal{A}\}})''_c \times \mathcal{A} \\ &= Y/_{\{(n,1),(1,n),\mathcal{B},\mathcal{A}\}} \times \mathcal{B} \times \mathcal{A} \stackrel{\text{sym}}{=} (Y''_c)''_r. \end{aligned}$$

Case II The RD-factor contains the CD-factor $(A \supset B)$, i.e. the RD-segment runs "on the inside" of the CD-segment (see figure 13). As in Case II, $(Y'_r)''_c = (Y''_c)'_r$. Note that $C_t \in A$, and therefore Y_A is a square. This justifies step (*) below, where A^{tr} denotes the top right cell of A. Note that A^{tr} has the same function in A as the cell (1, n) has in Y. The proof works even in the extreme case where $A^{tr} = (1, n)$.

$$\begin{aligned} (Y'_c)''_r &= (Y - C_t)''_r = (\mathcal{A} - C_t) \times Y_{\{(n,1),\mathcal{A}\}} \\ &= \mathcal{A}'_c \times Y_{\{(n,1),\mathcal{A}\}} = \left(\mathcal{A} \times Y_{\{(n,1),\mathcal{A}\}}\right)'_c = (Y''_r)'_c; \\ (Y''_r)''_c &= \left(\mathcal{A} \times Y_{\{(n,1),\mathcal{A}\}}\right)''_c = \mathcal{A}''_c \times Y_{\{(n,1),\mathcal{A}\}} \\ &= \mathcal{A}_{\{\mathcal{B},\mathcal{A}''\}} \times \mathcal{B} \times Y_{\{(n,1),\mathcal{A}\}}; \\ &\stackrel{(*)}{=} \left(Y_{\{(1,n),\mathcal{B}\}}\right)''_r \times \mathcal{B} = \left(Y_{\{(1,n),\mathcal{B}\}} \times \mathcal{B}\right)''_r = (Y''_c)''_r. \end{aligned}$$

Case III The CD-factor contains the RD-factor $(\mathcal{B} \supset \mathcal{A})$, i.e. the CD-segment runs "on the inside" of the RD-segment. This case is symmetric to Case II.

Case IV The RD- and CD-segments overlap (see figure 14). This happens exactly when the RD- and CD-segments differ from each other only in their final cells: the RD-segment intersects the rightmost column of *Y*, while the CD-segment intersects the bottom row of *Y*. Let $\mathcal{D} := \mathcal{B}/_{C_{b-1}} = \mathcal{A}/_{C_{t-1}}$. Since $\mathcal{B} - C_b$ and $\mathcal{A} - C_t$ contain exactly one bottom, resp. rightmost, cell, then $\mathcal{B} - C_b \equiv \mathcal{D} \equiv \mathcal{A} - C_t$ and $\mathcal{B}''_r = \mathcal{D} = \mathcal{A}'_c$. Clearly, $\mathcal{S} := Y/_{\mathcal{D}}$ is a *square*. Moreover, $\mathcal{S}_b = \mathcal{B}/_{\mathcal{D}} = C_{b-1}$, $\mathcal{S}_t = \mathcal{A}/_{\mathcal{D}} = C_{t-1}$; the top right cell of \mathcal{S} is the cell (1, n) of *Y*, and the

bottom left cell of S is the cell (n, 1) of Y. From this, we see that $Y_{\{\mathcal{C},(1,n)\}} = Y_{\{\mathcal{D},(n,1)\}}$ and $Y_{\{\mathcal{B},(1,n)\}} = Y_{\{\mathcal{A},(n,1)\}}$. Therefore,

$$\begin{aligned} (Y'_{r})''_{c} &= (Y - C_{b})''_{c} = \mathcal{B}/_{C_{b-1}} \times Y/_{\{\mathcal{B}/_{C_{b-1}}, (1,n)\}} \\ &= \mathcal{D} \times Y/_{\{\mathcal{D}, (1,n)\}} = \mathcal{D} \times Y/_{\{\mathcal{D}, (n,1)\}} \stackrel{\text{sym}}{=} (Y'_{c})''_{r}; \\ (Y''_{c})'_{r} &= (Y/_{\{\mathcal{B}, (1,n)\}} \times \mathcal{B})'_{r} = Y/_{\{\mathcal{B}, (1,n)\}} \times (\mathcal{B} - C_{b}) \\ &\equiv Y/_{\{\mathcal{B}, (1,n)\}} \times \mathcal{D} = Y/_{\{\mathcal{A}, (n,1)\}} \times \mathcal{D} \stackrel{\text{sym}}{=} (Y''_{r})'_{c}; \\ (Y''_{r})''_{c} &= (\mathcal{A} \times Y/_{\{\mathcal{A}, (n,1)\}})''_{c} = \mathcal{A}''_{c} \times Y/_{\{\mathcal{A}, (n,1)\}} \\ &= \mathcal{D} \times Y/_{\{\mathcal{A}, (n,1)\}} = \mathcal{D} \times Y/_{\{\mathcal{B}, (1,n)\}} \stackrel{\text{sym}}{=} (Y''_{c})''_{r}. \end{aligned}$$

The three special subcases when $C_{b-1} = (n, 1)$ (and hence $C_{t-1} = (1, n)$), when $C_b = (n, 1)$ (and hence $C_t = (1, n)$) and when $C_b = C_t$ (i.e. *Y* is a square), are easily checked to satisfy the desired equalities.

The discussion of these four cases completes the proof of Lemma 3.

We remark that, due to the degenerate nature of Case IV, two of the final cross-products turn out to be equal: $(Y_c'')_r' = (Y_c'')_r''$, and hence $(Y_c'')_r''$ has only two factors, rather than the three it has in Cases I–III.

5. New Wilf equivalences and consequences

Lemmas 2 and 3 imply the main result of the paper: (213) \sim^{s} (132), which we repeat below as Theorem 1. Combined with Proposition 1, this establishes a new class of Wilf-equivalent permutations.

Theorem 1 The permutations (213) and (132) are shape-Wilf-equivalent. Consequently, for any $\tau \in S_{n-3}$, the permutations $(n - 1, n - 2, n, \tau)$ and $(n - 2, n, n - 1, \tau)$ are Wilf-equivalent.

In particular, this completes the classification up to Wilf-equivalences of S_n , for $n \le 7$. Figure 5 lists the number of symmetry classes and Wilf-classes in each such S_n .

An amusing corollary about numerical equivalence of Young diagrams can be deduced from the above theorem and the row-decomposition formula. Recall that St is the standard staircase diagram. The *k*-staircase St_k is the Young diagram which consists of St plus the full k - 1 diagonals below the diagonal of St. In particular, St is St_1 , and the square $n \times n$ is

classes	S_1	S_2	S_3	S_4	S_5	S_6	S_7
symmetry	1	1	2	7	23	115	694
Wilf	1	1	1	3	16	91	595

Figure 5. Number of symmetry vs. Wilf classes in S_n , $n \le 7$.



Figure 6. $St_2 \cup St_2^1 \cup St_2^2 \cup St_2^3 \equiv St_2 \cup (St_2^2)^t \cup St_2^3 \cup St_2^1$.

 St_n . The *critical* staircase St_k of Y is the first staircase whose complement $Y \setminus St_k$ is a union of at least two connected components. Label such components by St_k^j , for j = 1, 2, ..., starting at the bottom left corner of Y. Thus, for every Y with critical staircase St_k , we have the *critical* decomposition $Y = St_k \cup_j St_k^j$.

Corollary 2 Let $Y = St_k \cup_{j=1}^{\ell} St_k^j$ be the critical decomposition of Y. The operations of permuting and of transposing the components St_k^j result in Young diagrams numerically equivalent to Y.

In other words, let $\tau \in S_{\ell}$ be any permutation, and let $\vec{t} = (t_1, t_2, \dots, t_{\ell}) \in \{1, t\}^{\ell}$ correspond to a choice $t_j = t$ to transpose, resp. $t_j = 1$ not to transpose, the component $St_k^{\tau(j)}$. Then the following (ordered) critical decompositions represent numerically equivalent Young diagrams (cf. figure 6):

$$St_k \bigcup_{j=1}^{\ell} St_k^j \equiv St_k \bigcup_{j=1}^{\ell} \left(St_k^{\tau(j)} \right)^{t_j}$$

Proof: We use induction on the size of *Y*. The initial cases are easily verified. Moreover, when there is only *one* "hanging" shape, the corollary simply states that transposing *Y* will yield a numerically equivalent diagram Y^t : this is the content of Theorem 1.

Suppose now that there are at least two hanging shapes: St_k^1 is the bottom shape, and let's name the remaining shapes the "upper" shapes. Choose a permutation σ and a transposition vector \vec{t} both of which leave the bottom shape St_k^1 fixed, and operate on the upper shapes of Y. Apply the row-decomposition formula to Y: when C runs over all bottom cells of Y, we have

$$Y \equiv \sum_{C} Y_{C} = \sum_{C} A_{C} \times B_{C}.$$
(11)

From the definition of the critical decomposition of Y, all RD-factors A_C are parts of the bottom shape St_k^1 , except for the first and the last RD-factors: there $A_C = Y_r$ or $A_C = \emptyset$. In any case, none of the upper shapes St_k^2 , St_k^3 ,... are broken up in (11). Thus, by induction hypothesis, we can apply σ and \vec{t} to each Y_C : $(Y_C^{\sigma})^{\vec{t}} = A_C \times (B_C^{\sigma})^{\vec{t}}$. We can then put back together the resulting diagrams via another application of the row-decomposition formula:

$$Y \equiv \sum_{C} Y_{C} = \sum_{C} A_{C} \times B_{C} \stackrel{\text{ind.}}{\equiv} \sum_{C} A_{C} \times \left(B_{C}^{\sigma}\right)^{\vec{t}} \equiv \left(Y^{\sigma}\right)^{\vec{t}}.$$

This shows that leaving the bottom shape fixed, we can permute and transpose the upper shapes in any way we like. We also know, again from Theorem 1, that transposing the whole diagram Y yields a numerically equivalent diagram Y^t . It is an easy exercise in algebra to verify that these two types of operations generate the whole group of operations required in Corollary 2.

The conclusion of Corollary 2 holds under a slightly relaxed hypothesis regarding which staircases can be used instead of the critical staircase: as long as the RD (or CD) formula breaks up only the bottom (or only the top) hanging shape, the above proof goes through without modifications. Further generalizations are also possible, for instance, applying recursively the Corollary just within a hanging shape. Finally, this can all be used to write down a generating function for the numbers $T_{(213)}(Y)$, but we will not do this here since it will take us too far a field.

6. Further discussion

A careful investigation of the new Wilf-pair (546213) ~ (465213) in S_6 leads to the observation that both permutation matrices can be decomposed into two blocks of 3 × 3 matrices, and further, that moving from one decomposition to the other involves a transposition of one of the blocks. Thus, one might be lead to conjecture that for any permutation matrices *A* and *B*, the following permutation matrices are Wilf-equivalent:

$$\left(\frac{A \mid 0}{0 \mid B}\right) \sim \left(\frac{A^t \mid 0}{0 \mid B}\right). \tag{12}$$

In order for (12) to give any new Wilf-classes, other than those obtained by symmetry or $I_t \stackrel{s}{\sim} J_t$, both *A* and *B* must be non-symmetric matrices. In *S*₆, there is only one such pair up to symmetry (denote the 1's by dots, and omit all 0's):

$$M(546213) = \left(\begin{array}{c} \cdot \\ \cdot \\ \hline \\ \cdot \\ \end{array}\right) \sim \left(\begin{array}{c} \cdot \\ \cdot \\ \hline \\ \cdot \\ \end{array}\right) = M(465213) \cdot$$

In S_7 , there are essentially 7 new Wilf-pairs which are covered by (12). Not surprisingly, these are the pairs appearing in figure 7. One possible approach to prove (12) would be to show $A \stackrel{s}{\sim} A^t$ for any permutation matrix A. Unfortunately, this is *not* true; it fails already in S_4 , e.g. (3142) $\stackrel{s}{\sim}$ (2413) since

$$\left|S_{(6,6,6,6,5,5)}(3142)\right| = 394 < 395 = \left|S_{(6,6,6,6,5,5)}(2413)\right|.$$

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$\tau \sim \sigma \text{ in } S_7$	$S_9(au)$	$S_{10}(\tau)$	$S_{11}(\tau)$	$S_{12}(\tau)$	$S_{13}(\tau)$
$(6571342) \sim (5761342)$	361300	3587768	38951398	457416920	5756026177
$(6571423) \sim (5761423)$	361300	3587768	38951411	457418106	5756088993
$(6572413) \sim (5762413)$	361300	3587768	38951430	457419793	5756176230
$(6572431) \sim (5762431)$	361300	3587768	38951467	457423216	5756360170
$(6574132) \sim (5764132)$	361300	3587780	38952330	457459680	5757549454
$(6571432) \sim (5761432)$	361301	3587834	38953996	457496956	5758168203
$(6571243) \sim (5761243)$	361301	3587834	38954024	457499462	5758298471

Figure 7. Final Wilf-equivalences in S_7 .

The Wilf-equivalence in (12) also *fails* in S_8 . For example:

$$M(68572413) = \left(\begin{array}{c} \cdot & \cdot \\ \cdot & \cdot \\ \hline & \cdot & \cdot \end{array}\right) \not\sim \left(\begin{array}{c} \cdot & \cdot \\ \cdot & \cdot \\ \hline & \cdot & \cdot \end{array}\right) = M(75862413) \cdot$$

The two permutations are equinumerant up to level 11, but they split on level 12:

 $|S_{12}(75862413)| = 476576750 < 476576751 = |S_{12}(68572413)|.$

This forces a reexamination of the 8 new Wilf-pairs in S_6 and S_7 . If we choose the "right" representatives of the symmetry classes, we can see that each permutation matrix contains a block corresponding to (213) or (132). This led to conjecturing and proving the shape-Wilf-equivalence (213) $\stackrel{s}{\sim}$ (132). As we can see, this SWE is far from coincidental, and it is the reason for the infinitely many new Wilf-equivalences of Theorem 1.

Let us now shift the emphasis of our discussion to a slightly different question. If we write down a table enumerating $|S_n(\tau)|$ for $\tau \in S_4$ as *n* increases, we notice a very plausible long-standing conjecture, which was first mentioned in [18] as a question:

Question 1 Is it true that if $|S_k(\tau)| < |S_k(\sigma)|$ for some k, then $|S_n(\tau)| < |S_n(\sigma)|$ for all $n \ge k$? In other words, modulo Wilf-equivalence, can we order linearly all permutations in S_n according to their relative restrictiveness: $\tau < \sigma$ if $|S_k(\tau)| < |S_k(\sigma)|$ for some k?

 S_4 is the first non-trivial case of Question 1. It was partially answered positively by Bóna in [4], where he shows $|S_n(1423)| < |S_n(1324)|$ for $n \ge 6$ and $|S_n(1234)| < |S_n(1324)|$ for $n \ge 7$, in relation to a conjecture of Wilf and Stanley. To the best of our knowledge, no one had published a counterexample to Question 1, until we found a counterexample in S_5 , followed by various types of counterexamples in S_6 and S_7 .

As figure 8 suggests, (53241) and (43251) *cannot* be ordered since $S_k(53241) < S_k(43251)$ for $k \le 12$, but $S_{13}(53241) > S_{13}(43251)$. Figures 9 and 10 list all counterexamples up to level 13 in S_6 , and some counterexamples in S_7 . The asterisks indicate the first level at which the corresponding permutations "switch" their relative restrictiveness, and hence cannot be ordered as in Conjecture 1. The "!!" in figure 9 refers to the permutation (546213), which is also part of the new Wilf-equivalence of Theorem 1.

$\tau \in S_5$	$S_7(\tau)$	$S_8(\tau)$	$S_9(au)$	$S_{10}(au)$	$S_{11}(\tau)$	$S_{12}(\tau)$	$S_{13}(\tau)$
(25314)	4578	33184	258757	2136978	18478134	165857600	1535336290
(31524)	4579	33216	259401	2147525	18632512	167969934	1563027614
(35214)	4579	33218	259483	2149558	18672277	168648090	1573625606
(35124)	4580	33249	260092	2159381	18815124	170605392	1599499163
(53124)	4580	33252	260202	2161837	18858720	171285237	1609282391
(42351)	4580	33252	260204	2161930	18861307	171341565	1610345257
(35241)	4580	33254	260285	2163930	18900534	172016256	1621031261
(53241)	4580	33256	260370	2166120	18945144	172810050	1633997788*
(43251)	4581	33283	260805	2171393	18994464	173094540	1632480259*
(32541)	4581	33284	260847	2172454	19015582	173461305	1638327423
(34215)	4581	33285	260886	2173374	19032746	173741467	1642533692
(31245)	4581	33286	260927	2174398	19053058	174094868	1648198050
(42315)	4581	33287	260967	2175379	19072271	174426353	1653484169
(12345)	4582	33324	261808	2190688	19318688	178108704	1705985883
(53421)	4582	33325	261853	2191902	19344408	178582940	1713999264
(52341)	4582	33325	261863	2192390	19358590	178904675	1720317763

Figure 8. Classification of S₅ up to Wilf equivalence.

$\tau \in S_7$	$S_9(\tau)$	$S_{10}(\tau)$	$S_{11}(\tau)$	$S_{12}(\tau)$	$S_{13}(\tau)$
253146	39424	344580	3283521	33633237*	366084190
523614	39425	344611	3283955	33632674*	365858205
352164	39425	344619	3284418	33648781*	366298292
426153	39425	344620	3284441	33648549*	366268369
365241	39425	344633	3285228	33676816*	367058930
236145	39426	344661	3285505	33671423*	366717782
356214	39426	344671	3286054	33689584	367192027*
315642	39426	344672	3286086	33689894	367181171*
315426	39426	344678	3286493	33705839*	367659357
524316	39426	344679	3286521	33705830*	367633865
325146	39426	344681	3286660	33711372*	367802282
532416	39426	344682	3286686	33711239*	367772447
326154	39427	344724	3287748	33732379*	368139497
156324	39427	344724	3287749	33732314*	368131670
643251	39427	344725	3287840	33736582*	368277055
453261	39427	344726	3287851	33735408*	368204877
435216	39427	344726	3287877	33737112	368270512*
532164	39427	344727	3287904	33737204	368255863*
632541	39427	344727	3287974*	33741840	368436494
513246	39427	344728	3287971*	33739711	368325636
624351	39427	344732	3288259	33751604	368699773*
652413	39427	344733	3288292	33751853	368681504*
643521	39427	344735	3288426	33757116	368841108*
546213	39428	344772	3289163	33765743	368833207*!!

Figure 9. Counterexamples to Question 1 in S_6 .

$\tau \in S_7$	$S_9(au)$	$S_{10}(au)$	$S_{11}(au)$	$S_{12}(au)$	$S_{13}(au)$
4725631	361297	3587548	38941876**	457096663*	5746716800
5427316	361297	3587549	38941876**	457093463*	5746500592
7536241	361297	3587572	38943767*	457184933	5749921907
2645173	361298	3587605	38943703*	457138045	5747366258

Figure 10. Some counterexamples to Question 1 in S_7 .

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$ au\in S_6$	$S_8(\tau)$	$S_9(au)$	$S_{10}(\tau)$	$S_{11}(au)$	$S_{12}(au)$	$S_{13}(\tau)$
(546213)	39428	344772	3289163	33765743	368833207	4247979687
(465213)	39428	344772	3289163	33765743	368833207	4247979687

Figure 11. Missing pair of Wilf-equivalence in S₆.



Figure 12. Case I for Y(9, 9, 9, 9, 8, 8, 5, 5, 5).

However, the above counterexamples do not preclude an *asymptotic* ordering of all permutations in S_n .

Definition 12 For $\tau, \sigma \in S_n$, we say that τ is *asymptotically smaller* than σ if $|S_k(\tau)| < |S_k(\sigma)|$ for all $k \gg 0$.

Conjecture 1 We can order asymptotically all permutations in S_n , modulo Wilfequivalence.

Regev [12] and Bóna [5] have worked on asymptotic behavior of certain types of permutations, but as of now, Conjecture 1 and some possible modifications of it are far from proven.

Finally, we haven't observed examples of permutations "switching" their relative positions more than once. Thus, a stronger version of the above conjecture is possible, where assymptotic ordering is replaced by the previous usual ordering, allowing for one switch: $\tau < _1\sigma$ if $|S_n(\tau) \ge |S_n(\sigma)|$ for $n \le K$ and $|S_n(\tau) \le |S_n(\sigma)|$ for n > K. Here K depends on τ and σ , and if K = 0, $\tau < \sigma$ in the strongest sense as in Question 1.



Figure 13. Case II for *Y*(9, 9, 9, 9, 9, 9, 7, 7, 4).



Figure 14. Case IV for *Y*(10, 10, 10, 9, 9, 8, 8, 6, 5, 3).

Conjecture 2 We can order linearly the Wilf-classes in S_n under the relation $<_1$.

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Notes

- 1. To keep the resemblance with the "shape" of τ , we coordinatize $M(\tau)$ from the bottom left corner.
- 2. In all tables, we skip the column corresponding to $|S_{n+1}(\tau)|$ if $\tau \in S_n$. It is easy to see that all permutations in S_n are equinumerant on level n + 1 (for example, cf. [11]).

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