Journal of Inequalities in Pure and Applied Mathematics

RATE OF CONVERGENCE OF THE DISCRETE POLYA ALGORITHM FROM CONVEX SETS. A PARTICULAR CASE

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©2000 Victoria University ISSN (electronic): 1443-5756 086-01



volume 3, issue 4, article 50, 2002.

Received 4 December, 2001; accepted 28 May, 2002.

Communicated by: A. Babenko



Abstract

In this work we deal with best approximation in ℓ_p^n , $1 , <math>n \ge 2$. For $1 , let <math>h_p$ denote the best ℓ_p^n -approximation to $f \in \mathbb{R}^n$ from a closed, convex subset K of \mathbb{R}^n , $f \notin K$, and let h^* be a best uniform approximation to f from K. In case that $h^* - f = (\rho_1, \rho_2, \cdots, \rho_n)$, $|\rho_j| = \rho$ for $j = 1, 2, \cdots, n$, we show that the behavior of $||h_p - h^*||$ as $p \to \infty$ depends on a property of separation of the set K from the ℓ_∞^n -ball $\{x \in \mathbb{R}^n : ||x - f|| \le \rho\}$ at $h^* - f$.

2000 Mathematics Subject Classification: 26D15

Key words: Best uniform approximation, Rate of convergence, Polya Algorithm, Strong uniqueness.

The authors gratefully acknowledge the many helpful suggestions of the referees during the revision of this paper.

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1. Introduction

Let (w_1, w_2, \ldots, w_n) be a fixed vector in \mathbb{R}^n , with $w_j > 0, j \in I_n := \{1, 2, \ldots, n\}$, $n \ge 2$. For $x = (x(1), x(2), \ldots, x(n)) \in \mathbb{R}^n$ we define

$$\|x\|_{p,w} := \left(\sum_{j=1}^{n} w_j |x(j)|^p\right)^{\frac{1}{p}}, \quad 1 \le p < \infty, \text{ and} \\ \|x\| := \max_{1 \le j \le n} |x(j)|.$$

Also we define $N := \sum_{j=1}^{n} w_j$.

Throughout the paper, K will always be a nonempty, closed, convex subset of \mathbb{R}^n . For $f \in \mathbb{R}^n \setminus K$, we will say that $h_{p,w} \in K$, $1 \leq p < \infty$, is a best $\ell_{n,w}^n$ -approximation to f from K if

$$||f - h_{p,w}||_{p,w} \le ||f - h||_{p,w} \quad \forall h \in K.$$

The existence of at least one best $\ell_{p,w}^n$ -approximation to f from K is a known fact for $1 \leq p < \infty$. Likewise, there always exists a best uniform approximation to f from K, i.e., an $h^* \in K$ that satisfies

 $||f - h^*|| \le ||f - h|| \quad \forall h \in K.$

We will henceforth assume f = 0 and $0 \notin K$. This causes no loss of generality, since all relevant properties are translation invariant. If $1 , there is a unique best <math>\ell_{p,w}^n$ -approximation. In this case, the next theorem [14] characterizes the best $\ell_{p,w}^n$ -approximation to 0 from K.



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Theorem 1.1 (Characterization of the best $\ell_{p,w}^n$ -**approximation).** Let K be a closed, convex subset of \mathbb{R}^n , $0 \notin K$. Then $h_{p,w}$, $1 \leq p < \infty$, is a best $\ell_{p,w}^n$ -approximation to 0 from K if and only if for all $h \in K$,

(1.1)
$$\sum_{j=1}^{n} w_j(h_{p,w}(j) - h(j)) |h_{p,w}(j)|^{p-1} \operatorname{sgn}(h_{p,w}(j)) \le 0, \quad \text{if } p > 1.$$

(1.2)
$$\sum_{j \in R(h_{1,w})} w_j \left(h_{1,w}(j) - h(j) \right) \operatorname{sgn}(h_{1,w}(j)) \le \sum_{j \in Z(h_{1,w})} w_j \left| h(j) \right|, \quad \text{if } p = 1,$$

where, if $g \in \mathbb{R}^n$, $Z(g) := \{j \in I_n : g(j) = 0\}$ and $R(g) := I_n \setminus Z(g)$.

It is also known [1, 6, 7] that if K is an affine subspace, then

(1.3)
$$\lim_{p \to \infty} h_{p,w} = h^*,$$

where in this case h^* is a particular best uniform approximation to 0 from K, called *strict* uniform approximation [12, 7] and whose definition is also valid in any closed, convex K. In [3, 8] it is proved that there exists a constant M > 0 such that $p ||h_{p,w} - h^*|| \le M$ for all p > 1. Moreover, from [13] it is deduced that there are constants $M_1, M_2 > 0$ and $0 \le a \le 1$, depending on K, such that

$$M_1 a^p \le p ||h_{p,w} - h^*|| \le M_2 a^p$$
 for all $p > 1$.

In [2, 7] it is shown that if K is not an affine subspace, then $h_{p,w}$ does not necessarily converge to the strict uniform approximation, though (1.3) is always



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valid whenever h^* is the unique best uniform approximation to 0 from K. In [6, 7] we can find sufficient conditions on K under which (1.3) is satisfied. In any case, the convergence of $h_{p,w}$ as $p \to \infty$ to a best uniform approximation is known as the **Polya algorithm** [11]. The purpose of this paper is to study the behavior of $||h_{p,w} - h^*||$ as $p \to \infty$ when h^* is a best uniform approximation to 0 from K and h^* satisfies $|h^*(j)| = \rho > 0 \forall j \in I_n$.





2. Relation Between Strong Uniqueness and Rate of Convergence

A useful concept in order to get a first general result on the rate of convergence of the Polya algorithm is strong uniqueness. It was established in 1963 by Newman and Shapiro [10] in the context of the uniform approximation to continuous functions by means of elements of a Haar space, although we could define it in any normed space.

Definition 2.1. Let $h^* \in K$ be a best uniform approximation to $0 \in \mathbb{R}^n$ from *K*. We say that h^* is strongly unique if there exists $\gamma > 0$ such that

(2.1)
$$||h - h^*|| \le \gamma(||h|| - ||h^*||) \quad \forall h \in K$$

It is obvious that if h^* is strongly unique, then h^* is the unique best uniform approximation to $0 \in \mathbb{R}^n$ from K.

Theorem 2.1. If the best uniform approximation h^* to 0 from K is strongly unique, then $p ||h_{p,w} - h^*||$ is bounded for all $p \ge 1$.

Proof. We first note that for every $h \in K$,

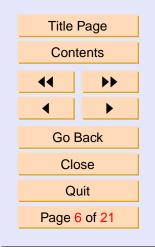
(2.2) $m^{\frac{1}{p}} \|h\| \le \|h\|_{p,w} \le N^{\frac{1}{p}} \|h\|,$

where $m := \min_{j \in I_n} \{w_j\}$. Let $\gamma > 0$ satisfy (2.1). Then for any $p \ge 1$,

(2.3)
$$||h_{p,w} - h^*|| \le \gamma(||h_{p,w}|| - ||h^*||)$$



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Applying (2.2) and the definition of best $\ell_{p,w}^n$ -approximation, we have

$$\begin{aligned} \|h_{p,w}\| - \|h^*\| &\leq \frac{1}{m^{\frac{1}{p}}} \|h_{p,w}\|_{p,w} - \|h^*\| \\ &\leq \frac{1}{m^{\frac{1}{p}}} \|h^*\|_{p,w} - \|h^*\| \\ &\leq \left[\left(\frac{N}{m}\right)^{\frac{1}{p}} - 1\right] \|h^*\| \\ &\leq \frac{(N-m)\|h^*\|}{mp}. \end{aligned}$$

From (2.3) we finally conclude that

$$p \|h_{p,w} - h^*\| \le \frac{\gamma(N-m)\|h^*\|}{m}$$
 for all $p \ge 1$.

The above inequality improves the proposal in [4] and [5].

2.1. The Particular Case $|h^*(j)| = \rho > 0, j = 1, 2, ..., n$

We henceforth suppose that $h^* \in K$ is a best uniform approximation to 0 from K, where $|h^*(j)| = \rho > 0$ for all $j \in I_n$. Under these conditions we will analyze the behaviour of $||h_{p,w} - h^*||$ as $p \to \infty$. In Theorem 2.3, our main result, we will prove that the converse of Theorem 2.1 – which is generally not true – is valid in this particular case. Since $\{x \in \mathbb{R}^n : ||x|| \le$



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 $\rho \} \cap K = \{h \in K : \|h\| = \rho \}, \text{ it is easy to see that there is a hyperplane} \\ \left\{ (x(1), x(2), \dots, x(n)) : \sum_{j=1}^{n} a_j \operatorname{sgn}(h^*(j)) x(j) = \rho \right\}, \text{ with } 0 \le a_j \le 1, \text{ all } \\ j \in I_n, \text{ and } \sum_{j=1}^{n} a_j = 1, \text{ that separates } K \text{ from the ball } \{x \in \mathbb{R}^n : \|x\| \le \rho \} \\ \text{ at } h^*, \text{ i.e., } \sum_{j=1}^{n} a_j \operatorname{sgn}(h^*(j)) h(j) \ge \rho \text{ for all } h \in K. \end{cases}$

Definition 2.2. We will say that

$$\pi := \left\{ (x(1), x(2), \dots, x(n)) : \sum_{j=1}^{n} a_j \operatorname{sgn}(h^*(j)) x(j) = \rho \right\}$$

is a hyperplane that strongly separates K from the ball $\{x \in \mathbb{R}^n : ||x|| \le \rho\}$ at h^* , or equivalently, that π is a strongly separating hyperplane at h^* , if

(2.4)
$$0 < a_j < 1, \ all \ j \in I_n, \quad \sum_{j=1}^n a_j = 1$$

and

(2.5)
$$\sum_{j=1}^{n} a_j \operatorname{sgn}(h^*(j)) h(j) \ge \rho \quad \forall h \in K.$$

In the proofs of Lemma 2.2 and Theorems 2.3 and 2.4 we will assume $h^*(j) = 1$ for all $j \in I_n$. This causes no loss of generality, since we can replace K by the closed, convex set

$$\left\{\tilde{h}\in\mathbb{R}^n:\tilde{h}(j)=\frac{1}{\rho}\,h(j)\,\mathrm{sgn}(h^*(j)),\ j\in I_n,\ h\in K\right\}.$$



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Lemma 2.2. If $p_k ||h_{p_k,w} - h^*||$ is bounded for $p_k \to \infty$, then there exists a strongly separating hyperplane at h^* .

Proof. Since $\lim_{p_k\to\infty} h_{p_k,w}(j) = h^*(j) = 1$, all $j \in I_n$, we can suppose $h_{p_k,w}(j) > 0$, all $j \in I_n$ and, without loss of generality, all p_k . Then, for every p_k the formula of characterization (1.1) can be expressed in the form

$$\sum_{j=1}^{n} w_j (h_{p_k, w}(j) - h(j)) h_{p_k, w}^{p_k - 1}(j) \le 0 \quad \forall h \in K.$$

Dividing by $||h_{p_k,w}||_{p_k,w}^{p_k}$, for every p_k we obtain

(2.6)
$$\sum_{j=1}^{n} w_j \left(\frac{h_{p_k,w}(j)}{\|h_{p_k,w}\|_{p_k,w}} \right)^{p_k} \frac{h(j)}{h_{p_k,w}(j)} \ge 1 \quad \forall h \in K.$$

Keeping in mind that

$$w_j h_{p_k,w}^{p_k}(j) \le \|h_{p_k,w}\|_{p_k,w}^{p_k} \le \|h^*\|_{p_k,w}^{p_k} = N, \quad j \in I_n$$

and after passage to a subsequence, we can suppose that $h_{p_k,w}^{p_k}(j)$, all $j \in I_n$, and $\|h_{p_k,w}\|_{p_k,w}^{p_k}$ are convergent. Now, by hypothesis, $p_k |h_{p_k,w}(j) - 1|$ is bounded for all $j \in I_n$ and all p_k . Hence we get

$$\lim_{p_k \to \infty} h_{p_k,w}^{p_k}(j) = \lim_{p_k \to \infty} \operatorname{Exp}\left(p_k(h_{p_k,w}(j) - 1)\right) > 0, \quad \text{all } j \in I_n.$$

Writing

$$a_j = \lim_{p_k \to \infty} w_j \left(\frac{h_{p_k, w}(j)}{\|h_{p_k, w}\|_{p_k, w}} \right)^{p_k}, \quad j \in I_n,$$



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we therefore deduce that $0 < a_j < 1$, all $j \in I_n$, and $\sum_{j=1}^n a_j = 1$. Taking limits as $p_k \to \infty$ in (2.6), we finally conclude that

$$\sum_{j=1}^{n} a_j h(j) \ge 1 \quad \forall h \in K.$$

Then $\left\{ (x(1), x(2), \dots, x(n)) := \sum_{j=1}^{n} a_j x(j) = 1 \right\}$ is a strongly separating hyperplane at h^* .

Theorem 2.3. The following statements are equivalent:

- (a) The best uniform approximation to 0 from K, h^* , is strongly unique.
- (b) $p \|h_{p,w} h^*\|$ is bounded for all $p \ge 1$.
- (c) $p_k ||h_{p_k,w} h^*||$ is bounded for a sequence $p_k \to \infty$.
- (d) There exists a strongly separating hyperplane at h^* .

Proof. (a) \Rightarrow (b) is Theorem 2.1. (b) \Rightarrow (c) is obvious. (c) \Rightarrow (d) is Lemma 2.2. To complete the theorem, we now prove (d) \Rightarrow (a). Suppose that there is a strongly separating hyperplane π at $h^* = (1, 1, ..., 1)$. Let $h \in K$. Observe that $||h|| \ge 1$. Let I_n^+ denote the subset of indices j in I_n such that $h(j) \ge 1$, and let $I_n^- := I_n \setminus I_n^+$. For all $j \in I_n^+$ we have $|h(j) - 1| = h(j) - 1 \le ||h|| - 1$. On the other hand, if $j \in I_n^-$, then |h(j) - 1| = 1 - h(j). Moreover, the inequality

$$\sum_{i \in I_n} a_i(h(i) - 1) \ge 0$$



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implies

$$a_{j}(1-h(j)) \leq \sum_{i \in I_{n}, i \neq j} a_{i}(h(i)-1)$$

$$\leq \sum_{i \in I_{n}^{+}} a_{i}(h(i)-1)$$

$$\leq (\|h\|-1) \sum_{i \in I_{n}^{+}} a_{i}$$

$$\leq (\|h\|-1)(1-a_{j}).$$

Thus, for all $j \in I_n$ we have

$$|h(j) - 1| \le \left(\frac{1}{\min_{i \in I_n} a_i} - 1\right) (||h|| - 1) := \gamma(||h|| - 1),$$

and so $||h - h^*|| \le \gamma(||h|| - ||h^*||)$.

Our goal now is to show that, under the conditions of Theorem 2.3, either $h_{p,w} = h^*$ for all p or there exist constants $M_1, M_2 > 0$ such that

$$M_1 \le p ||h_{p,w} - h^*|| \le M_2$$
 for all $p \ge 1$.

On the other hand, if there exists no strongly separating hyperplane at h^* , then the following example in \mathbb{R}^2 , where $\lim_{p\to\infty} h_{p,w} = h^*$, shows that the rate of convergence is as slow as we want.

Example 2.1. Let $\alpha : [1, +\infty) \to (0, 1]$ be a continuous strictly decreasing function such that $\alpha(1) = 1$ and $\lim_{t\to\infty} \alpha(t) = 0$ and let $\beta : (0, 1] \to [1, +\infty)$



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denote its inverse function, that will also be a strictly decreasing function. We define

$$f(x) := 1 + \int_{x}^{1} (1-t)^{\beta(1-t)} dt, \quad 0 \le x \le 1,$$

and let K be the convex hull of the set $\{(x, y) \in \mathbb{R}^2 : y = f(x), x \in [0, 1]\}$.

Observe that $h^* = (1, 1)$ is the unique best uniform approximation to (0, 0) from K. Moreover, the function f is smooth, convex and f'(1) = 0. This implies that the strongly separating hyperplane at h^* does not exist.

Let $h_p = (1 - \varepsilon_p, 1 + \delta_p)$ be the best *p*-approximation to (0, 0) from *K*, with $\varepsilon_p, \delta_p \downarrow 0$ as $p \to \infty$. Since the slopes of the curve y = f(x) and the ℓ_p -ball coincide at h_p , we have

$$\frac{(1-\varepsilon_p)^{p-1}}{(1+\delta_p)^{p-1}} = \varepsilon_p^{\beta(\varepsilon_p)}$$

and therefore

(2.7)
$$\lim_{p \to \infty} \varepsilon_p^{\beta(\varepsilon_p)/(p-1)} = \lim_{p \to \infty} \frac{1 - \varepsilon_p}{1 + \delta_p} = 1.$$

If $\varepsilon_p \leq \alpha(p)$, then $\beta(\varepsilon_p) \geq \beta(\alpha(p)) = p$, which contradicts (2.7). Then, for p large, we have $\varepsilon_p > \alpha(p)$. This shows that the rate of convergence of h_p to h^* as $p \to \infty$ can be as slow as we want.

Theorem 2.4. The following conditions are equivalent

(a) $h_{p,w} = h^*$ for all $p \ge 1$, (b) $h_{p_0,w} = h^*$ for some $p_0 \ge 1$,



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(c) the hyperplane

(2.8)
$$\pi := \left\{ (x(1), x(2), \dots, x(n)) : \sum_{j=1}^{n} \operatorname{sgn}(h^{*}(j)) \frac{w_{j}}{N} x(j) = \rho \right\}$$

is a strongly separating hyperplane at h^* .

Proof. (a) \Rightarrow (b) is obvious. (b) \Rightarrow (c) follows immediately from Theorem 1.1. Indeed, if $h_{p_0,w} = h^*$ for some $p_0 \ge 1$, then from (1.1) if $p_0 > 1$ or (1.2) if $p_0 = 1$, we have

(2.9)
$$\sum_{j=1}^{n} w_j(h^*(j) - h(j)) \operatorname{sgn}(h^*(j)) \le 0 \quad \forall h \in K,$$

which is equivalent to the fact that π is a strongly separating hyperplane at h^* . Also from (1.1) and (1.2), the inequality (2.9) implies that $h_{p,w} = h^*$ for all $p \ge 1$ and so (c) \Rightarrow (a).

Theorem 2.5. Suppose that $h_{p,w} \neq h^*$ for some $p \geq 1$ and there exists a strongly separating hyperplane at h^* . Then there are constants $M_1, M_2 > 0$ such that

$$M_1 \le p ||h_{p,w} - h^*|| \le M_2 \text{ for all } p \ge 1.$$

Proof. Assume that there exists a strongly separating hyperplane at h^* , where $h^*(j) = 1$ for all $j \in I_n$. From Theorem 2.3, there is a constant $M_2 > 0$ such that

$$p ||h_{p,w} - h^*|| \le M_2.$$



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Therefore, to prove the theorem it is sufficient to show that $\inf_{p\geq 1}\{p \| h_{p,w} - h^* \| \} > 0$. Suppose the contrary. In order to get a contradiction, we only need to consider the two following exhaustive cases:

1. There exists a sequence $p_k \to \infty$ such that $\lim_{p_k \to \infty} p_k ||h_{p_k} - h^*|| = 0$. In this case $\lim_{p_k \to \infty} p_k |h_{p_k,w}(j) - 1| = 0$ for all $j \in I_n$. This implies that $h_{p_k,w}^{p_k}(j) \to 1$ as $p_k \to \infty$ and

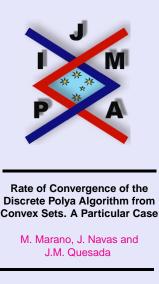
$$a_j^* = \lim_{p_k \to \infty} w_j \left(\frac{h_{p_k, w}(j)}{\|h_{p_k, w}\|_{p_k, w}} \right)^{p_k} = \frac{w_j}{N}, \qquad j = 1, 2, \dots, n,$$

which means (see the proof of Lemma 2.2) that the hyperplane (2.8), with $h^*(j) = \rho = 1$ for all $j \in I_n$, is a strongly separating hyperplane at h^* . From Theorem 2.4 (c), $h_{p,w} = h^*$ for all $p \ge 1$, which contradicts the hypothesis of the theorem.

2. There exists a sequence $p_k \to p_0$, $1 \le p_0 < \infty$, such that $\lim_{p_k \to p_0} p_k ||h_{p_k,w} - h^*|| = 0$. Since $h_{p_k,w} \to h_{p_0,w}$, we deduce that $||h_{p_0,w} - h^*|| = \lim_{p_k \to p_0} ||h_{p_k,w} - h^*|| = 0$ and so $h_{p_0,w} = h^*$. Now, using the statement (b) of Theorem 2.4, we conclude that $h_{p,w} = h^*$, for all $p \ge 1$. A contradiction.

2.2. A Numerical Example in Isotonic Approximation

Let
$$f = (\underbrace{a+1, \ldots, a+1}_{r}, \underbrace{a-1, \ldots, a-1}_{n-r}) \in \mathbb{R}^{n}$$
, and let K be the convex set





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of the nondecreasing vectors in \mathbb{R}^n , i.e.

$$K = \{ h \in \mathbb{R}^n : h(i) \le h(j) \ \forall i, j \in I_n, i < j \}$$

In this case, the (unique) best uniform approximation to f from K is the element $h^* = (a, a, \ldots, a)$. Thus $h_{p,w} \to h^*$ as $p \to \infty$. Furthermore, it is easy to see that

$$h_{p,w} = (x_{p,w}, x_{p,w}, \dots, x_{p,w}) \in \mathbb{R}^n, \quad 1$$

for some $x_{p,w}$ satisfying $a - 1 \le x_{p,w} \le a + 1$.

In order to translate h^* to a vertex of the ℓ_∞^n -ball, we consider the closed, convex set

$$\widetilde{K} = \{ \widetilde{h} \in \mathbb{R}^n : \widetilde{h}(j) = h(j) - f(j), \ j \in I_n, \ h \in K \}.$$

In this way we obtain

•
$$\widetilde{f} = (0, 0, \dots, 0);$$

• $\widetilde{h}^* = h^* - f = (\underbrace{-1, \dots, -1}_{r}, \underbrace{1, \dots, 1}_{n-r})$

To simplify the notation, we will write $\sigma_j = \text{sgn}(\tilde{h}^*(j)), j \in I_n$. Now, we are interested in obtaining a strongly separating hyperplane at \tilde{h}^* , i.e., a hyperplane

$$\pi := \left\{ (x(1), x(2), \dots, x(n)) : \sum_{j=1}^{n} a_j \, \sigma_j \, x(j) = 1 \right\}$$

such that



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(p1)
$$0 < a_j < 1$$
, all $j \in I_n$, and $\sum_{1}^{n} a_j = 1$;
(p2) $\sum_{j=1}^{n} \sigma_j a_j \widetilde{h}(j) \ge 1 \forall \widetilde{h} \in \widetilde{K}$.

Proposition 2.6. Let $S := \sum_{j=1}^{r} w_j$. Then the above hyperplane π , with

$$a_j = \frac{w_j}{2S}$$
 if $1 \le j \le r$, $a_j = \frac{w_j}{2(N-S)}$ if $r+1 \le j \le n$,

satisfies (p1) and (p2), and therefore it is a strongly separating hyperplane at \tilde{h}^* .

Proof. By definition, $0 < a_j < 1$ for all $j \in I_n$. Furthermore,

$$\sum_{j=1}^{n} \sigma_j a_j \widetilde{h}^*(j) = \sum_{j=1}^{n} a_j = \sum_{j=1}^{r} \frac{w_j}{2S} + \sum_{j=r+1}^{n} \frac{w_j}{2(N-S)} = \frac{1}{2} + \frac{1}{2} = 1.$$

Then (p1) holds.

Since

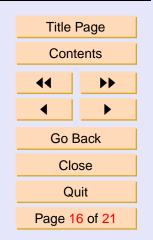
$$\sum_{j=1}^{n} \sigma_j a_j f(j) = -(a+1) \sum_{j=1}^{r} \frac{w_j}{2S} + (a-1) \sum_{j=r+1}^{n} \frac{w_j}{2(N-S)} = -1,$$

(p2) is equivalent to

(2.10) $\sum_{j=1}^{n} \sigma_j a_j h(j) \ge 0 \quad \forall h \in K.$



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But if h is a nondecreasing vector, then (2.10) is immediate because

$$\sum_{j=1}^{r} w_j h(j) \le h(r) \sum_{j=1}^{r} w_j = S h(r)$$

and

$$\sum_{j=r+1}^{n} h(j) \ge h(r) \sum_{j=r+1}^{n} w_j = (N-S)h(r),$$

and therefore

$$\sum_{j=1}^{n} \sigma_j a_j h(j) = -\frac{1}{2S} \sum_{j=1}^{r} w_j h(j) + \frac{1}{2(N-S)} \sum_{j=r+1}^{n} w_j h(j)$$
$$\geq -\frac{1}{2S} S h(r) + \frac{1}{2(N-S)} (N-S) h(r) = 0.$$

This concludes the proof.

From Proposition 2.6 we deduce that if S = N/2, then

(2.11)
$$\left\{ (x(1), x(2), \dots, x(n)) : \sum_{j=1}^{n} \sigma_j \, \frac{w_j}{N} \, x(j) = 1 \right\}$$

is a strongly separating hyperplane at \tilde{h}^* , and from Theorem 2.4 this is equivalent to $\tilde{h}_{p,w} = \tilde{h}^*$ for all $1 \le p < \infty$. In the case that $S \ne N/2$, we claim that $\tilde{h}_{p,w} \to \tilde{h}^*$ as $p \to \infty$ exactly at a rate $O\left(\frac{1}{p}\right)$. From Proposition 2.6 and Theorems 2.4 and 2.5 we only need to show that (2.11) is not a strongly separating



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hyperplane at \tilde{h}^* . This last assertion is true since (2.10), with $a_j = w_j/N$, all $j \in I_n$, implies

(2.12)
$$\sum_{j=r+1}^{n} w_j h(j) \ge \sum_{j=1}^{r} w_j h(j) \quad \forall h \in K.$$

On the other hand, if S < N/2 then $h = (-1, -1, ..., -1) \in K$ does not satisfy (2.12), and an analogous conclusion is valid for $h = (1, 1, ..., 1) \in K$ if S > N/2. This proves the claim.

In what follows we obtain these same results calculating directly the best $\ell_{p,w}^n$ -approximations to f from K, namely, $h_{p,w} = (x_{p,w}, x_{p,w}, \dots, x_{p,w})$. It is easy to check that

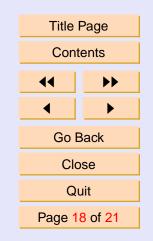
$$x_{p,w} = \frac{a - 1 + a\left(\frac{S}{N-S}\right)^{\frac{1}{p}} + \left(\frac{S}{N-S}\right)^{\frac{1}{p}}}{1 + \left(\frac{S}{N-S}\right)^{\frac{1}{p}}}, \quad 1$$

Then we immediately conclude that if S = N/2, then $h_{p,w} = h^*$ for p > 1, and if $S \neq N/2$, then $h_{p,w} \rightarrow h^*$ as $p \rightarrow \infty$. Moreover, we can calculate the rate of convergence. Indeed,

$$\lim_{p \to \infty} \frac{h_{p,w}(j) - h^*(j)}{1/p} = \lim_{p \to \infty} \frac{\frac{\left(\frac{S}{N-S}\right)^{1/p} - 1}{1 + \left(\frac{S}{N-S}\right)^{1/p}}}{1/p}$$
$$= \frac{1}{2} \lim_{p \to \infty} \frac{\left(\frac{S}{N-S}\right)^{\frac{1}{p}} - \frac{1}{N}}{1/p}$$



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$$=\frac{1}{2}\ln\left(\frac{S}{N-S}\right).$$

The rate of convergence is exactly $O\left(\frac{1}{p}\right)$.



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