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SOME RESULTS ON A GENERALIZED USEFUL INFORMATION MEASURE

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Abstract

A parametric mean length is defined as the quantity

$$\label{eq:Lu} \begin{split} {}_{\alpha\beta}L_u &= \frac{\alpha}{\alpha-1} \left[1 - \sum P_i^\beta \left(\frac{u_i}{\sum u_i p_i^\beta} \right)^{\frac{1}{\alpha}} D^{-n_i(\frac{\alpha-1}{\alpha})} \right], \\ & \text{where } \alpha \neq 1, \ \sum p_i = 1 \end{split}$$

this being the useful mean length of code words weighted by utilities, u_i . Lower and upper bounds for $_{\alpha\beta}L_u$ are derived in terms of useful information for the incomplete power distribution, p^{β} .

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1. Introduction

Consider the following model for a random experiment S,

$$S_N = [E; P; U]$$

where $E = (E_1, E_2, \ldots, E_n)$ is a finite system of events happening with respective probabilities $P = (p_1, p_2, \ldots, p_N)$, $p_i \ge 0$, and $\sum p_i = 1$ and credited with utilities $U = (u_1, u_2, \ldots, u_N)$, $u_i > 0$, $i = 1, 2, \ldots, N$. Denote the model by E, where

(1.1)
$$E = \begin{bmatrix} E_1 & E_2 & \cdots & E_N \\ p_1 & p_2 & \cdots & p_N \\ u_1 & u_2 & \cdots & u_N \end{bmatrix}$$

We call (1.1) a Utility Information Scheme (UIS). Belis and Guiasu [3] proposed a measure of information called 'useful information' for this scheme, given by

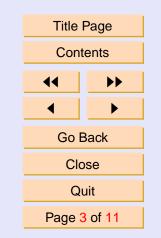
(1.2)
$$H(U;P) = -\sum u_i p_i \log p_i,$$

where H(U; P) reduces to Shannon's [8] entropy when the utility aspect of the scheme is ignored i.e., when $u_i = 1$ for each *i*. Throughout the paper, \sum will stand for $\sum_{i=1}^{N}$ unless otherwise stated and logarithms are taken to base D (D > 1).

Guiasu and Picard [5] considered the problem of encoding the outcomes in (1.1) by means of a prefix code with codewords w_1, w_2, \ldots, w_N having lengths



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 n_1, n_2, \ldots, n_N and satisfying Kraft's inequality [4]

(1.3)
$$\sum_{i=1}^{N} D^{-n_i} \le 1,$$

where D is the size of the code alphabet. The useful mean length L_u of code was defined as

(1.4)
$$L_u = \frac{\sum u_i n_i p_i}{\sum u_i p_i}$$

and the authors obtained bounds for it in terms of H(U; P).

Longo [8], Gurdial and Pessoa [6], Khan and Autar [7], Autar and Khan [2] have studied generalized coding theorems by considering different generalized measures of (1.2) and (1.4) under condition (1.3) of unique decipherability.

In this paper, we study some coding theorems by considering a new function depending on the parameters α and β and a utility function. Our motivation for studying this new function is that it generalizes some entropy functions already existing in the literature (see C. Arndt [1]). The function under study is closely related to Tsallis entropy which is used in physics.



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2. Coding Theorems

Consider a function

(2.1)
$$_{\alpha\beta}H(U;P) = \frac{\alpha}{\alpha-1} \left[1 - \left(\frac{\sum u_i p_i^{\alpha\beta}}{\sum u_i p_i^{\beta}}\right)^{\frac{1}{\alpha}} \right],$$

where $\alpha > 0 \ (\neq 1), \beta > 0, p_i \ge 0, i = 1, 2, ..., N$ and $\sum p_i \le 1$.

- (i) When $\beta = 1$ and $\alpha \to 1$, (2.1) reduces to a measure of useful information for the incomplete distribution due to Belis and Guiasu [3].
- (ii) When $u_i = 1$ for each *i* i.e., when the utility aspect is ignored, $\sum p_i = 1$, $\beta = 1$ and $\alpha \to 1$, the measure (2.1) reduces to Shannon's entropy [10].
- (iii) When $u_i = 1$ for each *i*, the measure (2.1) becomes entropy for the β -power distribution derived from *P* studied by Roy [9]. We call $_{\alpha\beta}H(U; P)$ in (2.1) the generalized useful measure of information for the incomplete power distribution P^{β} .

Further consider,

(2.2)
$$\alpha_{\beta}L_{u} = \frac{\alpha}{\alpha - 1} \left[1 - \sum P_{i}^{\beta} \left(\frac{u_{i}}{\sum u_{i} p_{i}^{\beta}} \right)^{\frac{1}{\alpha}} D^{-n_{i}(\frac{\alpha - 1}{\alpha})} \right],$$

where $\alpha > 0 \ (\neq 1), \sum p_i \leq 1$.



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- (i) For $\beta = 1$, $u_i = 1$ for each i and $\alpha \to 1$, $_{\alpha\beta}L_u$ in (2.2) reduces to the useful mean length L_u of the code given in (1.4).
- (ii) For $\beta = 1$, $u_i = 1$ for each i and $\alpha \to 1$, $_{\alpha\beta}L_u$ becomes the optimal code length defined by Shannon [10].

We establish a result, that in a sense, provides a characterization of $_{\alpha\beta}H(U;P)$ under the condition of unique decipherability.

Theorem 2.1. For all integers D > 1

(2.3)
$$_{\alpha\beta}L_u \ge {}_{\alpha\beta}H(U;P)$$

under the condition (1.3). Equality holds if and only if

(2.4)
$$n_i = -\log\left(\frac{u_i P_i^{\alpha\beta}}{\sum u_i p_i^{\alpha\beta}}\right).$$

Proof. We use Hölder's [11] inequality

(2.5)
$$\sum x_i y_i \ge \left(\sum x_i^p\right)^{\frac{1}{p}} \left(\sum y_i^q\right)^{\frac{1}{q}}$$

for all $x_i \ge 0$, $y_i \ge 0$, i = 1, 2, ..., N when $P < 1 \ (\ne 1)$ and $p^{-1} + q^{-1} = 1$, with equality if and only if there exists a positive number c such that

Setting

$$x_i = p_i^{\frac{\alpha\beta}{\alpha-1}} \left(\frac{u_i}{\sum u_i p_i^{\beta}}\right)^{\frac{1}{\alpha-1}} D^{-n_i},$$



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$$y_i = p_i^{\frac{\alpha\beta}{1-\alpha}} \left(\frac{u_i}{\sum u_i p_i^\beta}\right)^{\frac{1}{1-\alpha}},$$

 $p = 1 - 1/\alpha$ and $q = 1 - \alpha$ in (2.5) and using (1.3) we obtain the result (2.3) after simplification for $\frac{\alpha}{\alpha - 1} > 0$ as $\alpha > 1$.

Theorem 2.2. For every code with lengths $\{n_i\}$, i = 1, 2, ..., N, $_{\alpha\beta}L_u$ can be made to satisfy,

(2.7)
$$\alpha_{\beta}L_{u} \geq \alpha_{\beta}H(U;P)D^{(\frac{1-\alpha}{\alpha})} + \frac{\alpha}{1-\alpha}\left[1-D^{(\frac{1-\alpha}{\alpha})}\right].$$

Proof. Let n_i be the positive integer satisfying, the inequality

(2.8)
$$-\log\left(\frac{u_i P_i^{\alpha\beta}}{\sum u_i p_i^{\alpha\beta}}\right) \le n_i < -\log\left(\frac{u_i P_i^{\alpha\beta}}{\sum u_i p_i^{\alpha\beta}}\right) + 1.$$

Consider the intervals

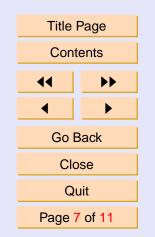
(2.9)
$$\delta_i = \left[-\log\left(\frac{u_i P_i^{\alpha\beta}}{\sum u_i p_i^{\alpha\beta}}\right), -\log\left(\frac{u_i P_i^{\alpha\beta}}{\sum u_i p_i^{\alpha\beta}}\right) + 1 \right]$$

of length 1. In every δ_i , there lies exactly one positive number n_i such that

(2.10)
$$0 < -\log\left(\frac{u_i P_i^{\alpha\beta}}{\sum u_i p_i^{\alpha\beta}}\right) \le n_i < -\log\left(\frac{u_i P_i^{\alpha\beta}}{\sum u_i p_i^{\alpha\beta}}\right) + 1.$$



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It can be shown that the sequence $\{n_i\}$, i = 1, 2, ..., N thus defined, satisfies (1.3). From (2.10) we have

(2.11)
$$n_{i} < -\log\left(\frac{u_{i}P_{i}^{\alpha\beta}}{\sum u_{i}p_{i}^{\alpha\beta}}\right) + 1$$
$$\Rightarrow D^{-n_{i}} < \left(\frac{u_{i}P_{i}^{\alpha\beta}}{\sum u_{i}p_{i}^{\alpha\beta}}\right)D$$
$$\Rightarrow D^{-n_{i}(\frac{\alpha-1}{\alpha})} < \left(\frac{u_{i}P_{i}^{\alpha\beta}}{\sum u_{i}p_{i}^{\alpha\beta}}\right)^{\frac{1-\alpha}{\alpha}}D^{\frac{\alpha-1}{\alpha}}$$

Multiplying both sides of (2.11) by $p_i^{\beta} \left(\frac{u_i}{\sum u_i p_i^{\alpha\beta}}\right)^{\frac{1}{\alpha}}$, summing over i = 1, 2, ..., N and simplifying, gives (2.7).

Theorem 2.3. For every code with lengths $\{n_i\}$, i = 1, 2, ..., N, of Theorem 2.1, $_{\alpha\beta}L_u$ can be made to satisfy

(2.12)
$$_{\alpha\beta}H(U;P) \leq _{\alpha\beta}L_u < _{\alpha\beta}H(U;P) + \frac{\alpha}{\alpha-1}(1-D)$$

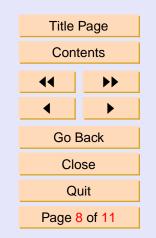
Proof. Suppose

(2.13)
$$\overline{n_i} = -\log\left(\frac{u_i P_i^{\alpha\beta}}{\sum u_i p_i^{\alpha\beta}}\right)$$

Clearly $\overline{n_i}$ and $\overline{n_i} + 1$ satisfy 'equality' in Hölder's inequality (2.5). Moreover, $\overline{n_i}$ satisfies Kraft's inequality (1.3).



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Suppose n_i is the unique integer between $\overline{n_i}$ and $\overline{n_i} + 1$, then obviously, n_i satisfies (1.3).

Since $\alpha > 0 \ (\neq 1)$, we have

$$(2.14) \qquad \sum p_{i}^{\beta} \left(\frac{u_{i}}{\sum u_{i}p_{i}^{\beta}}\right)^{\frac{1}{\alpha}} D^{\overline{n}_{i}(\alpha-1)/\alpha}$$

$$\leq \sum p_{i}^{\beta} \left(\frac{u_{i}}{\sum u_{i}p_{i}^{\beta}}\right)^{\frac{1}{\alpha}} D^{n_{i}(\alpha-1)/\alpha}$$

$$< D \left(\sum p_{i}^{\beta} \left(\frac{u_{i}}{\sum u_{i}p_{i}^{\beta}}\right)^{\frac{1}{\alpha}} D^{\overline{n}_{i}(\alpha-1)/\alpha}\right)$$

Since,

$$\sum p_i^{\beta} \left(\frac{u_i}{\sum u_i p_i^{\beta}} \right)^{\frac{1}{\alpha}} D^{\overline{n}_i(\alpha-1)/\alpha} = \left(\frac{\sum u_i p_i^{\alpha\beta}}{\sum u_i p_i^{\beta}} \right)^{\frac{1}{\alpha}}$$

Hence, (2.14) becomes

$$\left(\frac{\sum u_i p_i^{\alpha\beta}}{\sum u_i p_i^{\beta}}\right)^{\frac{1}{\alpha}} \le \sum p_i^{\beta} \left(\frac{u_i}{\sum u_i p_i^{\beta}}\right)^{\frac{1}{\alpha}} D^{-\overline{n}_i(\alpha-1)/\alpha} < D \left(\frac{\sum u_i p_i^{\alpha\beta}}{\sum u_i p_i^{\beta}}\right)^{\frac{1}{\alpha}}$$

which gives the result (2.12).



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